

Category of elements

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Definition 1. Given a locally small category \mathcal{A} and a functor $X : \mathcal{A}^{op} \rightarrow \mathbf{Set}$, the **category of elements** of X , denoted $\mathbb{E}(X)$ or $\int^{\mathcal{A}} X$, is defined as follows:

- Objects are pairs $(A \in \mathbb{A}, x \in X(A))$,
- Morphisms $f : (A, x) \rightarrow (A', x')$ are maps $f : A \rightarrow A' \in \mathcal{A}$ such that $(Xf)(x') = x$.

Given a presheaf X , there is a projection functor $P : \mathbb{E}(X) \rightarrow \mathcal{A}$ that sends $(A, x) \mapsto A$ and $f \mapsto f$. As a result of the property that morphisms satisfy, we can write them as $f : (A', (Xf)(x)) \rightarrow (A, x)$. It is worth noticing that if there is an \mathbb{A} -morphism $f : A' \rightarrow A$, then there is a unique element $x' \in X(A')$ such that there is an $\mathbb{E}(X)$ -morphism $f : (A', x') \rightarrow (A, x)$, namely $x' = (Xf)(x)$. The category of elements can also be treated as a comma category.

Lemma 2. [2, Exercise 6.2.22] *There is an isomorphism $\mathbb{E}(X) \cong (1 \rightrightarrows X)$.*

Proof. We look at the comma category for the following diagram:

$$\begin{array}{ccc} & \mathbb{A}^{op} & \\ & \downarrow X & \\ \mathbb{1} & \xrightarrow{1} & \mathbf{Set} \end{array}$$

Here, $\mathbb{1}$ is the terminal category and $1 : \mathbb{1} \rightarrow \mathbf{Set}$ is the functor that selects the terminal set. This category has as objects, pairs $(A \in \mathbb{A}, x : 1 \rightarrow X(A))$ and morphisms $f : (A, x) \rightarrow (A', x')$ are commuting triangles:

$$\begin{array}{ccc} \mathbb{1} & \xrightarrow{x'} & X(A') \\ & \searrow x & \downarrow Xf \\ & & X(A) \end{array}$$

That this triangle commutes is the same as stating $x = (Xf)(x')$, which is the condition above. \square

Proposition 3. [2, Exercise 6.2.23] *Let X be a presheaf on a locally small category. X is representable if and only if $\mathbb{E}(X)$ has a terminal object.*

Proof. The category $\mathbb{E}(X)$ has a terminal object if and only if there is an object (A, x) such that for any (A', x') , there is exactly one morphism $f : (A', x') \rightarrow (A, x)$. This is equivalent to there being an $A \in \mathcal{A}$ and $x \in X(A)$ such that for all $A' \in \mathcal{A}$, $x \in X(A')$, there is a unique morphism $f : A' \rightarrow A$ such that $(Xf)(x) = x'$. This condition is equivalent to X being representable, by [2, Corollary 4.3.2]. \square

The category of elements is very useful when looking at colimits in functor categories.

Proposition 4. [2, Theorem 6.2.17] *Let \mathbb{A} be small and $X : \mathbb{A}^{op} \rightarrow \mathbf{Set}$ a presheaf. Then X is the colimit of the following diagram:*

$$\mathbb{E}(X) \xrightarrow{P} \mathcal{A} \xrightarrow{H_\bullet} [\mathcal{A}^{op}, \mathbf{Set}]$$

That is, $X \cong \lim_{\rightarrow \mathbb{E}(X)} (H_\bullet \circ P)$.

We should first note that this does make sense; as \mathcal{A} is small, so is $\mathbb{E}(X)$, hence a colimit does indeed exist.

Proof. We know that presheaf categories have all (small) limits and colimits, so a colimit of $H_\bullet \circ P$ exists. Let $Y \in [\mathbb{A}^{op}, \mathbf{Set}]$ be a presheaf and let $(\alpha_{(A,x)} : (H_\bullet \circ P)(A, x) \rightarrow Y)_{(A,x) \in \mathbb{E}(X)}$ be a cocone on $H_\bullet \circ P$ with vertex Y . We can simply this to have $(\alpha_{(A,x)} : (H_A \rightarrow Y)_{(A,x) \in \mathbb{E}(X)})$. This is a family of natural transformations, so for all $f : (A', x') \rightarrow (A, x)$ in $\mathbb{E}(X)$, the following diagram commutes

$$\begin{array}{ccc} H_{A'} & & \\ H_f \downarrow & \searrow^{\alpha_{(A', (Xf)(x))}} & \\ H_A & \xrightarrow{\alpha_{(A,x)}} & Y \end{array} \quad (1)$$

By the Yoneda lemma, every natural transformation $\alpha_{(A,x)} : H_A \rightarrow Y$ corresponds to a unique element $(\alpha_{(A,x)})_A(1_A) \in Y(A)$, which we shall denote $y_{(A,x)}$. As diagram (1) commutes, it commutes for all $A \in \mathbb{A}$, so in particular it commutes for A' . This gives us the following:

$$\begin{array}{ccc} H_{A'}(A') & & 1_{A'} \longmapsto (\alpha_{(A', (Xf)(x))}_{A'}(1_{A'}) \quad (= y_{(A', (Xf)(x))}) \\ H_f(A') \downarrow & \searrow^{\alpha_{(A', (Xf)(x))}_{A'}} & \downarrow \\ H_A(A') \xrightarrow{(\alpha_{(A,x)})_{A'}} Y(A') & & f \longmapsto (\alpha_{(A,x)})_{A'}(f) \end{array} \quad (2)$$

This gives us $y_{(A', (Xf)(x))} = (\alpha_{(A,x)})_{A'}(f)$. As $\alpha_{(A,x)}$ is a natural transformation, the following square commutes:

$$\begin{array}{ccc}
H_A(A) & \xrightarrow{H_A(f)} & H_A(A') \\
(\alpha_{(A,x)})_A \downarrow & & \downarrow (\alpha_{(A,x)})_{A'} \\
Y(A) & \xrightarrow{Yf} & Y(A')
\end{array}
\quad
\begin{array}{ccc}
1_A & \xrightarrow{\quad\quad\quad} & f \\
\downarrow & & \downarrow \\
(\alpha_{(A,x)})_A(1_A) & \longmapsto & (Yf)((\alpha_{(A,x)})_A(1_A))
\end{array}
\tag{3}$$

This gives us $(Yf)(y_{(A,x)}) = (\alpha_{(A,x)})_{A'}(f)$. Combining this with the above we see that a cocone on Y is a collection of elements $(y_{(A,x)})_{(A,x) \in \mathbb{E}(X)}$ such that for any $f : (A', (Xf)(x)) \rightarrow (A, x)$ in $\mathbb{E}(X)$, $(Yf)(y_{(A,x)}) = y_{(A', (Xf)(x))}$.

An equivalent way to write $y_{(A,x)}$ is $\bar{\alpha}_A(x) : X(A) \rightarrow Y(A)$ and treat it as a function. The properties above then say for any $f : (A', (Xf)(x)) \rightarrow (A, x)$ in $\mathbb{E}(X)$, $(Yf)(\bar{\alpha}_A(x)) = \bar{\alpha}_{A'}((Xf)(x))$, that is to say the following diagram commutes for all f :

$$\begin{array}{ccc}
X(A) & \xrightarrow{Xf} & X(A') \\
\bar{\alpha}_A \downarrow & & \downarrow \bar{\alpha}_{A'} \\
Y(A) & \xrightarrow{Yf} & Y(A')
\end{array}
\tag{4}$$

This shows that $\bar{\alpha} : X \rightarrow Y$ is a natural transformation. As all of the above is equivalent, we see that a cocone on Y is the same as a map from X into Y , hence X is the colimit of $H_\bullet \circ P$. We can write this as equivalence formally as

$$[\mathbb{E}(X), [\mathbb{A}^{op}, \mathbf{Set}]](H_\bullet \circ P, \Delta Y) \cong [\mathbb{A}^{op}, \mathbf{Set}](X, Y).$$

This is an application of the dual of [2, Equation 6.2]. \square

0.1 An equivalence

Given a set S , there is an equivalence of categories $\mathbf{Set}/S \simeq \mathbf{Set}^S$, where the latter has as objects S indexed tuples of sets. Given $(A, f : A \rightarrow S) \in \mathbf{Set}/S$, we form the tuple $(f^{-1}(s))_{s \in S}$ and given a tuple $(A_s)_{s \in S}$, we form the disjoint union $\coprod_{s \in S} A_s$ along with the function $g : \coprod_{s \in S} A_s \rightarrow S$ that sends every element in each A_s to s . This equivalence can be abstracted to categories by the following theorem.

Theorem 5. [1, Proposition 1.1.7] *Let \mathbb{A} be a small category and $X : \mathbb{A}^{op} \rightarrow \mathbf{Set}$ a presheaf on \mathbb{A} . Then there is an equivalence of categories*

$$[\mathbb{A}^{op}, \mathbf{Set}]/X \simeq [\mathbb{E}(X)^{op}, \mathbf{Set}]. \tag{5}$$

Proof. There are lots of naturality conditions that need to be checked; however, we shall ignore most of them as they are quite easy to check. We first define the following functor:

$$\begin{aligned} \widehat{\cdot} : [\mathbb{A}^{op}, \mathbf{Set}]/X &\rightarrow [\mathbb{E}(X)^{op}, \mathbf{Set}] \\ (F, \alpha : F \rightarrow X) &\mapsto ((\widehat{F}, \widehat{\alpha}) : \mathbb{E}(X)^{op} \rightarrow \mathbf{Set}), \\ (\lambda : (F, \alpha) \rightarrow (G, \beta)) &\mapsto (\widehat{\lambda} : (\widehat{F}, \widehat{\alpha}) \rightarrow (\widehat{G}, \widehat{\beta})). \end{aligned}$$

The functor $(\widehat{F}, \widehat{\alpha})$ is defined as follows:

$$\begin{aligned} (\widehat{F}, \widehat{\alpha}) : \mathbb{E}(X)^{op} &\rightarrow \mathbf{Set} \\ (A, x) &\mapsto \alpha_A^{-1}(x), \\ f : (A', (Xf)(x)) &\rightarrow (A, x) \mapsto (\widehat{F}, \widehat{\alpha})(f) : \alpha_A^{-1}(x) \rightarrow \alpha_{A'}^{-1}((Xf)(x)). \end{aligned}$$

Where $(\widehat{F}, \widehat{\alpha})(f)(y) = (Ff)(y)$. The natural transformation $\widehat{\lambda}$ has components $\widehat{\lambda}_{(A,x)} : \alpha_A^{-1} \rightarrow \beta_A^{-1}(x)$ with $\widehat{\lambda}_{(A,x)}(y) = \lambda_A(y)$. We now define a map in the other direction:

$$\begin{aligned} \widetilde{\cdot} : [\mathbb{E}(X)^{op}, \mathbf{Set}] &\rightarrow [\mathbb{A}^{op}, \mathbf{Set}]/X \\ P : \mathbb{E}(X)^{op} \rightarrow \mathbf{Set} &\mapsto \left(\widetilde{P}_A : \coprod_{x \in X(A)} P_x(A) \rightarrow X(A) \right)_{A \in \mathbb{A}}, \\ \lambda : P \rightarrow Q &\mapsto \left(\widetilde{\lambda}_A : \coprod_{x \in X(A)} P_x(A) \rightarrow \coprod_{x \in X(A)} Q_x(A) \right)_{A \in \mathbb{A}}. \end{aligned}$$

The functor $P_x : \mathbb{A}^{op} \rightarrow \mathbf{Set}$ is defined as $P_x(A) = P(A, x)$. This can then be made into a functor $\coprod_{x \in X(-)} P_x : \mathbb{A}^{op} \rightarrow \mathbf{Set}$. The natural transformation \widetilde{P} has components defined by the universal property of the coproduct. If $y \in P_x(A)$ then $\widetilde{P}_A(y) = x$. The natural transformation $\widetilde{\lambda}$ has components with the following action on $y \in P(A, x)$ - $\widetilde{\lambda}_A(y) = \lambda_{(A,x)}(y)$.

We need to show that the composites of these functors are naturally isomorphic to the identity functors. Given $(F, \alpha) \in [\mathbb{A}^{op}, \mathbf{Set}]/X$, $(\widehat{F}, \widehat{\alpha})$ is a pair $(\coprod_{x \in X(-)} \alpha_{(-)}^{-1}(x), \widehat{\alpha})$. For any $A \in \mathbb{A}$, there is a map $\varphi_A^{(F, \alpha)}$ such that the following commutes:

$$\begin{array}{ccc} \coprod_{x \in X(A)} \alpha_A^{-1}(x) & \xrightarrow{\widehat{\alpha}_A} & X(A) \\ \varphi_A^{(F, \alpha)} \downarrow & \nearrow \alpha_A & \\ F(A) & & \end{array} \quad \begin{array}{ccc} y & \longrightarrow & \alpha_A(y) \\ \downarrow & \nearrow & \\ y & & \end{array}$$

This makes sense as for every $x \in X(A)$, $\alpha_A^{-1}(x) \subseteq F(A)$, so $\varphi_A^{(F,\alpha)}(y)$ is an inclusion map. We notice that each $\alpha_A^{-1}(x)$ is disjoint, so $\varphi_A^{(F,\alpha)}$ is injective. It must also be surjective as for any $y \in F(A)$, $y \in \alpha_A^{-1}(\alpha_A(y))$, hence $\varphi_A^{(F,\alpha)}$ is an isomorphism for all A . We need first show that it is natural in $A \in \mathbb{A}$, so let $f : A' \rightarrow A$ be an \mathbb{A} -morphism. We see that the following commutes:

$$\begin{array}{ccc}
\widetilde{(F,\alpha)}(A) & \xrightarrow{\widetilde{(F,\alpha)}(f)} & \widetilde{(F,\alpha)}(A') \\
\varphi_A^{(F,\alpha)} \downarrow & & \downarrow \varphi_{A'}^{(F,\alpha)} \\
F(A) & \xrightarrow{F(f)} & F(A')
\end{array}
\quad
\begin{array}{ccc}
\coprod_{x \in X(A)} \alpha_A^{-1}(x) & \xrightarrow{\widetilde{(F,\alpha)}(f)} & \coprod_{x \in X(A')} \alpha_{A'}^{-1}(x) \\
\varphi_A^{(F,\alpha)} \downarrow & & \downarrow \varphi_{A'}^{(F,\alpha)} \\
F(A) & \xrightarrow{F(f)} & F(A')
\end{array}$$

$$\begin{array}{ccc}
y & \longrightarrow & (Ff)(y) \\
\downarrow & & \downarrow \\
y & \longrightarrow & (Ff)(y)
\end{array}$$

Now we have a morphism $\varphi^{(F,\alpha)} : \widetilde{(F,\alpha)} \rightarrow (F,\alpha)$, for every $(F,\alpha) \in [\mathbb{A}^{op}, \mathbf{Set}]$. We need to show that it is natural. Let $\lambda : (F,\alpha) \rightarrow (G,\beta)$ be a $[\mathbb{A}^{op}, \mathbf{Set}]/X$ -morphism. It suffices to show that the following diagram commutes for every $A \in \mathbb{A}$, which we see it does.

$$\begin{array}{ccc}
\widetilde{(F,\alpha)}(A) & \xrightarrow{\tilde{\lambda}_A} & \widetilde{(G,\beta)}(A) \\
\varphi_A^{(F,\alpha)} \downarrow & & \downarrow \varphi_A^{(G,\beta)} \\
F(A) & \xrightarrow{\lambda_A} & G(A)
\end{array}
\quad
\begin{array}{ccc}
\coprod_{x \in X(A)} \alpha_A^{-1}(x) & \xrightarrow{\tilde{\lambda}_A} & \coprod_{x \in X(A)} \beta_A^{-1}(x) \\
\varphi_A^{(F,\alpha)} \downarrow & & \downarrow \varphi_A^{(G,\beta)} \\
F(A) & \xrightarrow{\lambda_A} & F(A)
\end{array}$$

$$\begin{array}{ccc}
y \in \alpha_A^{-1}(x) & \longrightarrow & \lambda_A(y) \\
\downarrow & & \downarrow \\
y & \longrightarrow & \lambda_A(y)
\end{array}$$

This shows that φ is natural, hence $\tilde{\cdot} \cong \mathbb{1}_{[\mathbb{A}^{op}, \mathbf{Set}]/X}$, naturally. The other isomorphism is shown similarly. \square

References

- [1] T. LEINSTER, *Higher operads, higher categories*, 2003.

- [2] ———, *Basic category theory*. Cambridge Studies in Advanced Mathematics, Vol. 143, Cambridge University Press, 2014, 2016. Accessed: 25/06/18.